

# Leveraging augmented-Lagrangian techniques for differentiating over infeasible quadratic programs in machine learning

Antoine Bambade<sup>1,2</sup>

<sup>1</sup>*Inria and ENS Paris : Willow and Sierra teams*

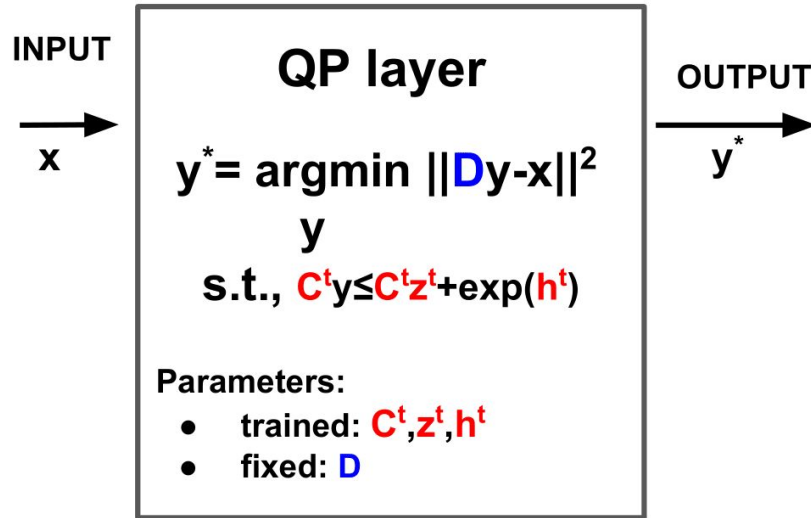
<sup>2</sup>*École des Ponts Paris Tech*

Joint work with Fabian Schramm, Adrien Taylor,  
Justin Carpentier



# Quadratic programming layer pipeline

More recent literature considers differentiable optimization problems as layers.



**Figure:** Example of a Quadratic Programming Layer (with  $\mathbf{D}$  nonsingular)

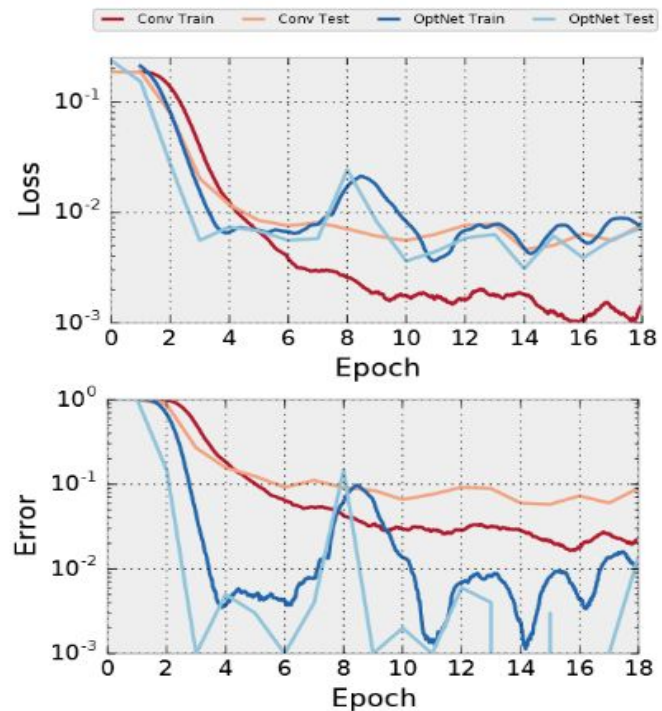
# QP layers in machine learning

Convex QP layers performs better than a ConvNet for solving Sudokus.

			3
1			
		4	
4			1

2	4	1	3
1	3	2	4
3	1	4	2
4	2	3	1

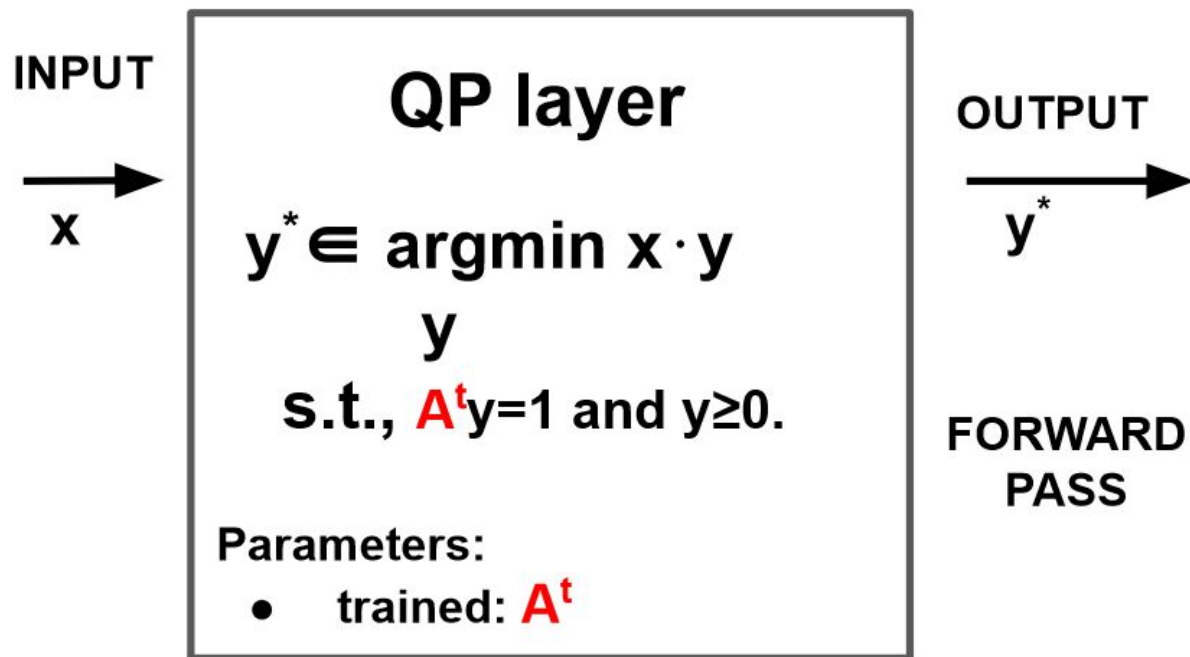
**Figure:** Example of Sudoku.



**Figure:** Training and test plots<sup>1</sup>.

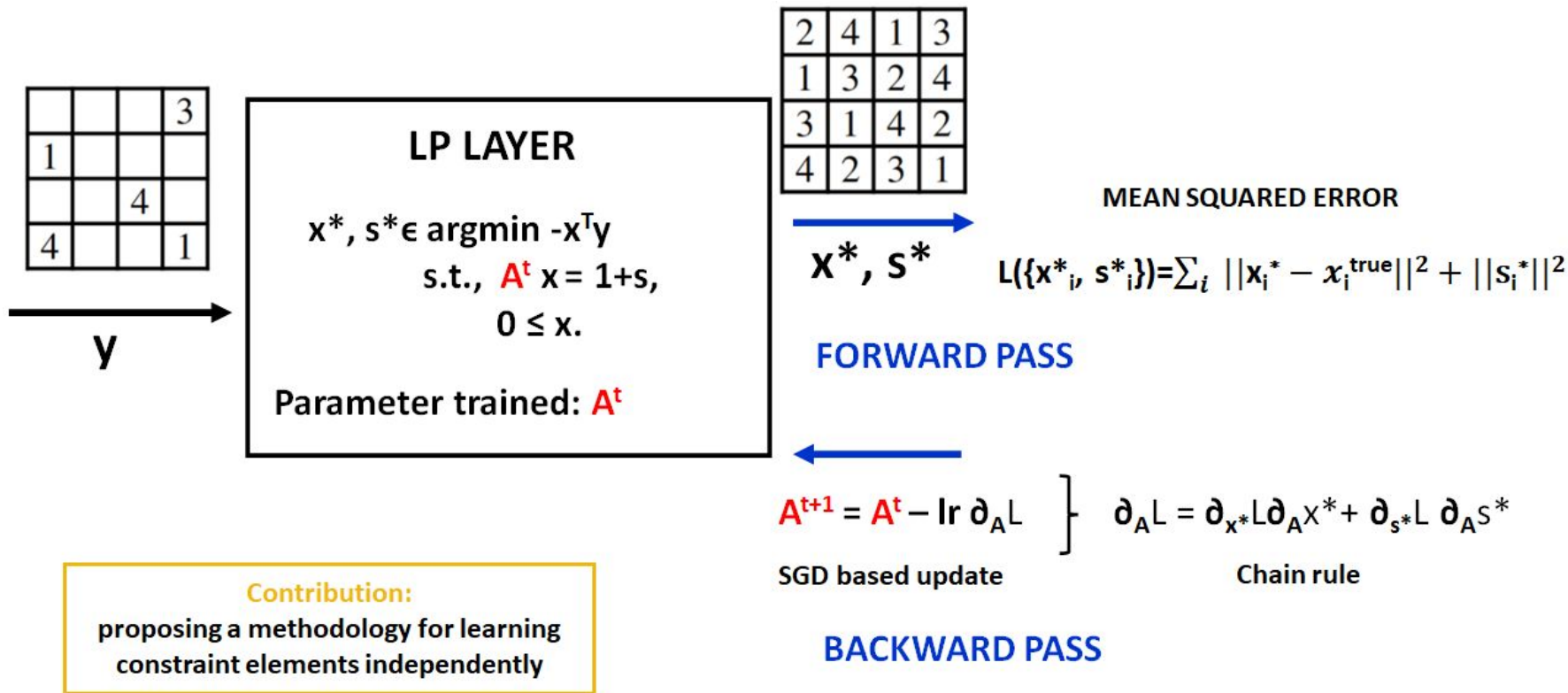
<sup>1</sup>B. Amos, Z. Kolter (2021)

## QP layers cons: limited trainable architecture



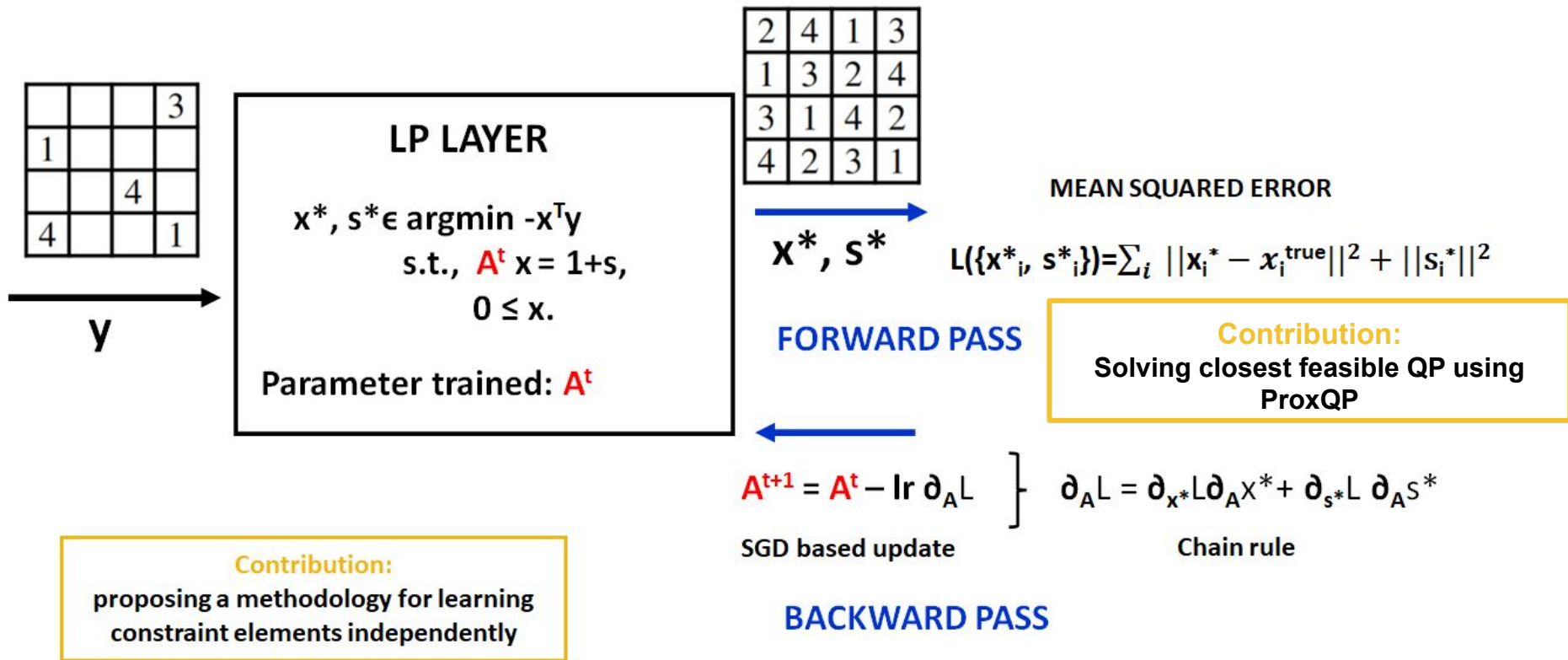
**Figure:** a LP layer. Nothing guarantees during training that the vector of 1 lies in the range space of  $A^t$ .

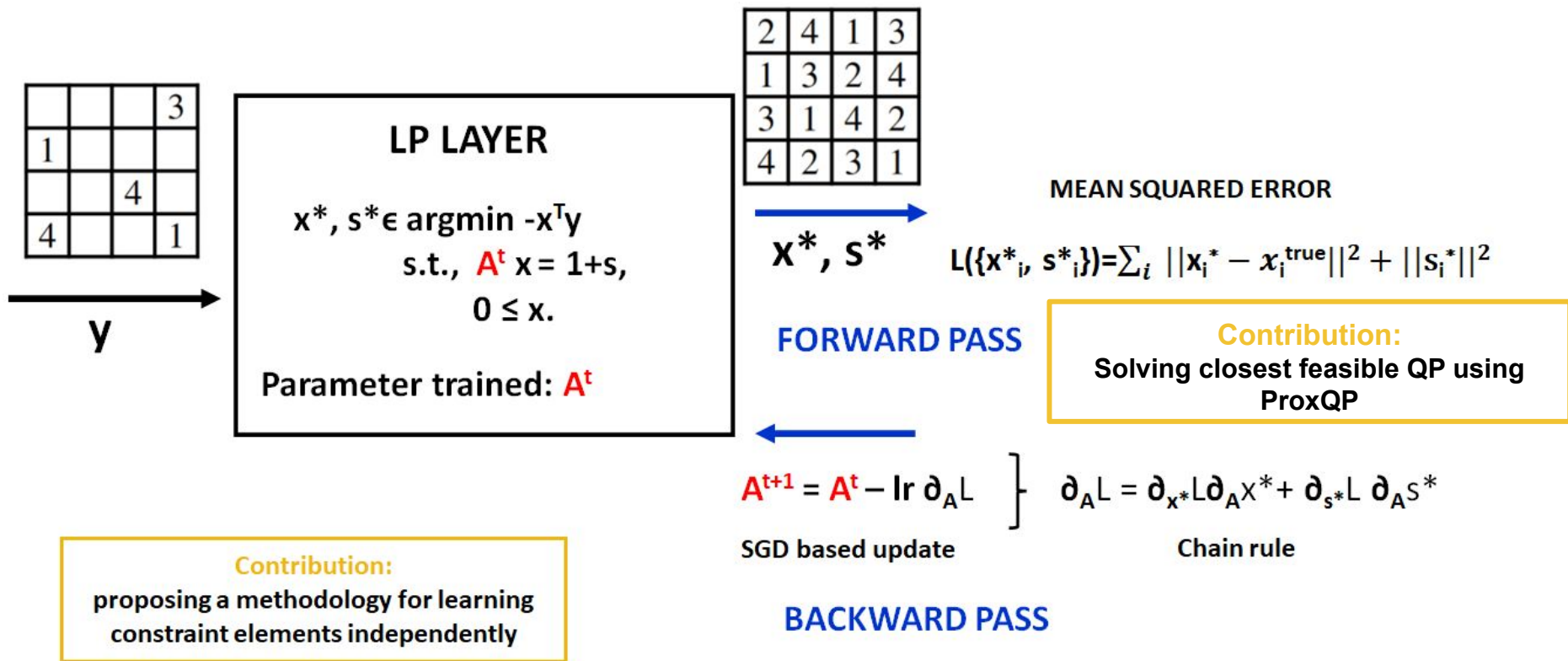
# Solution outline: ideal pipeline



**Contribution:**  
 proposing a methodology for learning  
 constraint elements independently

# Solution outline: ideal pipeline





			3
1			
		4	
4			1

$y$

**LP LAYER**

$x^*, s^* \in \operatorname{argmin} -x^T y$   
s.t.,  $A^t x = 1 + s,$   
 $0 \leq x.$

Parameter trained:  $A^t$

2	4	1	3
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$x^*, s^*$

**FORWARD PASS**



$A^{t+1} = A^t - lr \partial_A L$   
SGD based update

**BACKWARD PASS**

**Contribution:**  
**QPLayer: A full differentiable pipeline in C++ connected with PyTorch.**

MEAN SQUARED ERROR

$$L(\{x_i^*, s_i^*\}) = \sum_i ||x_i^* - x_i^{true}||^2 + ||s_i^*||^2$$

**Contribution:**  
**solving closest feasible QP using ProxQP**

$$\partial_A L = \partial_{x^*} L \partial_A x^* + \partial_{s^*} L \partial_A s^*$$

Chain rule

**Contribution:**  
**proposing a methodology for learning constraint elements independently**

**Contribution:**  
**propose algorithms to differentiate through closest feasible QP solutions**



## Software contribution

# ProxSuite

THE ADVANCED PROXIMAL OPTIMIZATION TOOLBOX

License **BSD 2-Clause** docs **online** CI - Linux/OSX/Windows - Cond **passing** pypi package **0.6.1** Anaconda.org **0.6.1**

- ✓ **fast:** C++ implementation, with homemade linear Cholesky solver
- ✓ **scalable:** various backends for dense, sparse and matrix-free optimization
- ✓ **easy-to-use:** API closed to OSQP, Python and Julia bindings
- ✓ **open-source:** BSD-license, easily installable

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</> Development: <https://github.com/simple-robotics/proxsuite>

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## Summary

PyPI link

<https://pypi.org/project/proxsuite>

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A classical technique: **the Implicit Function Theorem.**

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$$s^*(\theta) = \arg \min_{s \in \mathbb{R}^{n_i}} \frac{1}{2} \|s\|_2^2$$
$$\text{s.t. } x^*(\theta), z^*(\theta) \in \arg \min_{x \in \mathbb{R}^n} \max_{z \in \mathbb{R}_+^{n_i}} L(x, z, s; \theta),$$

with  $L(x, z, s; \theta) \triangleq f(x; \theta) + z^\top (C(\theta)x - u(\theta) - s)$ .

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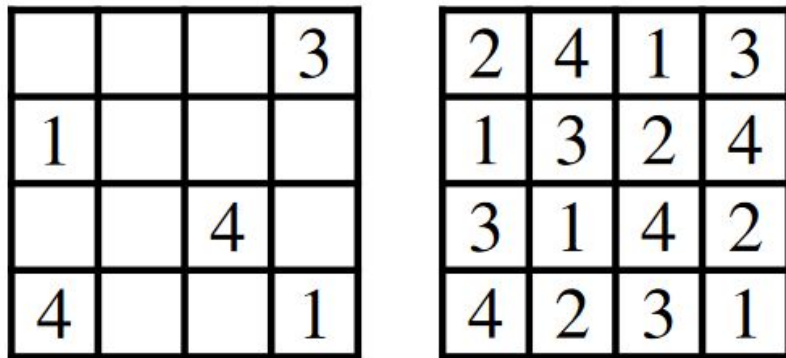
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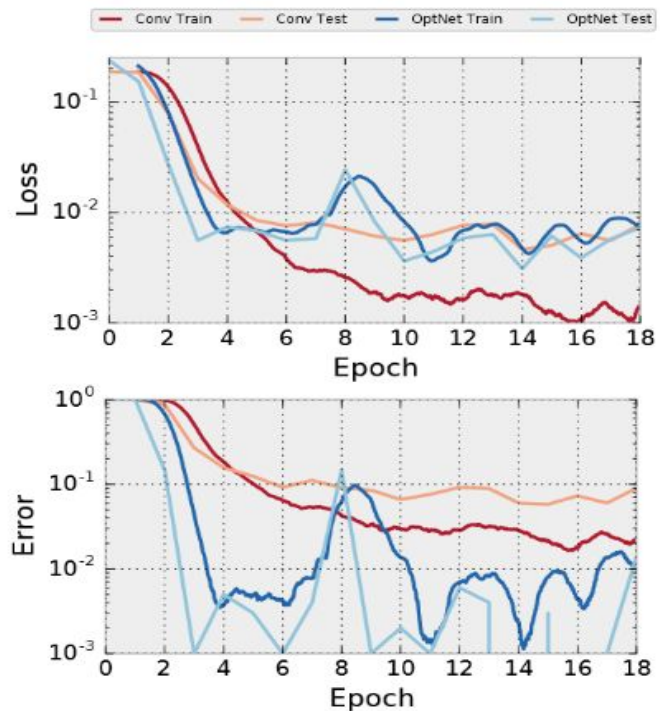
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# Numerical benchmark: back to the Sudoku problem.

Convex QP layers performs better than a ConvNet for solving Sudokus.



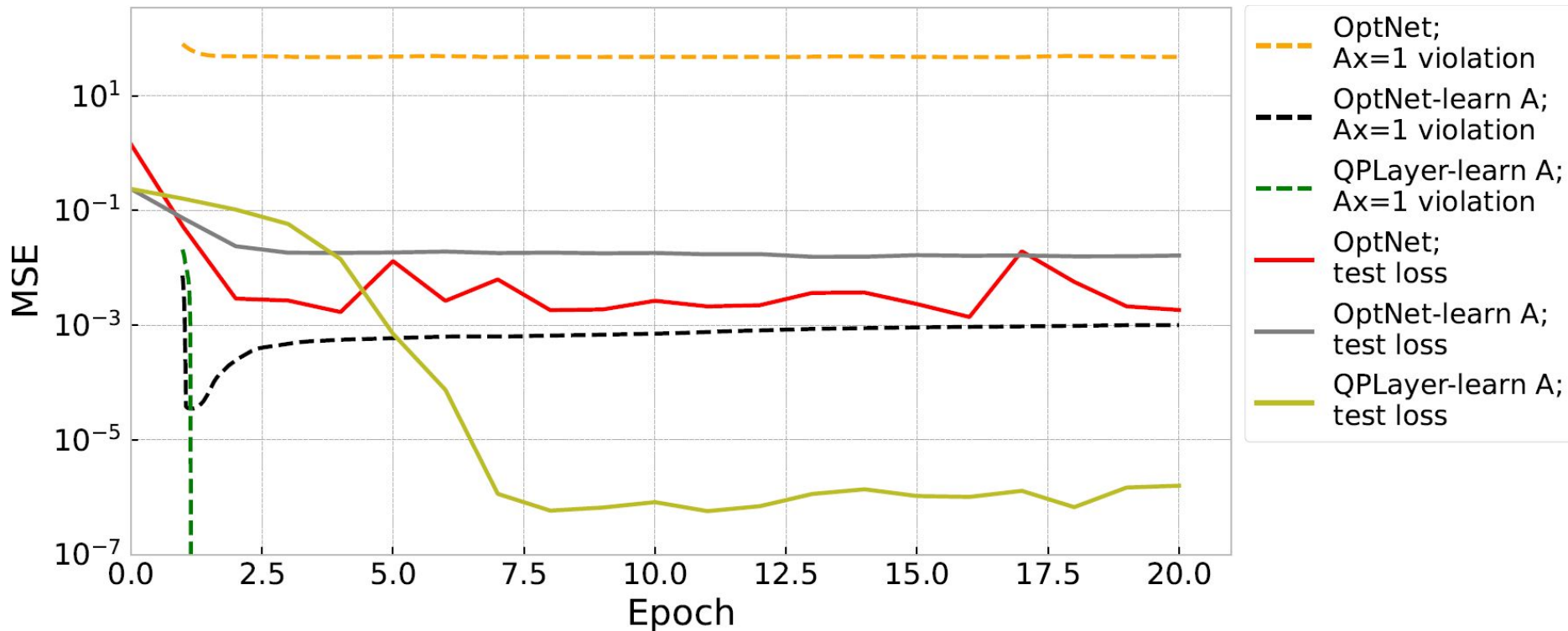
**Figure:** Example of Sudoku.



**Figure:** Training and test plots<sup>1</sup>.

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# Loss comparison



- **Methodology for learning new QP layers**
  - IFT for closest feasible QPs
  - Extended conservative Jacobians
- **QPlayer: open-source differentiable pipeline**
  - Use Augmented-Lagrangian techniques
  - Connected with PyTorch

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<https://github.com/Simple-Robotics/proxsuite>



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**Software contribution**

Exact Method of Multipliers launched from  $z^0$   
 with  $\sum_k 1/\mu^k$  finite

$$z^{k+1} = z^k - 1/\mu^k s^k, \quad s^k \in \partial\delta(z^{k+1})$$

